

# Data Lab @ UOI

## Members:



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...+ 14  
students

## Research Topics:

- Data Warehousing (ETL and OLAP)
- Data Visualization
- Graph Data Analytics, Evolving Graphs
- Query Results Diversification
- Social Media Data Analysis
- Spatial Data Management and Analysis

# D.A.T.A.

**Data Algorithms Technologies Architectures**



# University of Ioannina



# Outline

- CineCubes: Aiding data workers gain insights from OLAP queries (Panos Vassiliadis, [pvassil@cs.uoi.gr](mailto:pvassil@cs.uoi.gr))
- Analysis of large social networks that evolve over time (Evaggelia Pitoura, [pitoura@cs.uoi.gr](mailto:pitoura@cs.uoi.gr))
- Extracting Knowledge from Social Networks and Their Content (Panayiotis Tsaparas, [tsap@cs.uoi.gr](mailto:tsap@cs.uoi.gr))
- Search and Analysis of spatially-enriched data (Nikos Mamoulis, [nikos@cs.uoi.gr](mailto:nikos@cs.uoi.gr))

# Caught somewhere in time



- Query result = (just) a set of tuples
- No difference from the 70's when this assumption was established and tailored for
  - what people had available then
    - ... a green/orange monochrome screen
    - ... a dot-matrix(?) printer
    - ... nothing else
  - users being programmers



NO MORE JUST SETS OF TUPLES!

REPLACE QUERY ANSWERING WITH  
INSIGHT GAINING!

... but how?



# Cinecubes produce **small stories** presented as **data movies** ...

to orthogonally combine the following tasks:

- STORY* • expand a query result with the results of **complementary queries** which allow the user to **contextualize and analyze** the information content of the original query.
- with  
INSIGHTS* • extract meaningful, important patterns, or “**highlights**” from the query results
- Presented  
as a  
MOVIE* • present the results (a) **properly visualized**; (b) **enriched with** an automatically extracted **text** that comments on the result; (c) **vocally enriched**, i.e., enriched with audio that allows the user not only to see, but also hear

<http://www.cs.uoi.gr/~pvassil/projects/cinecubes/>

1		Assoc	Post-grad	Some-college	University
	Gov	40.73	43.58	38.38	42.14
	Private	41.06	45.19	38.73	43.06
	Self-emp	46.68	47.24	45.70	46.61

### Original query

Here, you can see the answer of the original query. You have specified education to be equal to 'Post-Secondary', and work at level 1. We report on Avg of Hrs grouped by education at level 2, and work at level 1. We highlight the largest values with red and the lowest values with blue.

Column Some-college has 2 of the 3 lowest values.  
Row Self-emp has 3 of the 3 highest values.  
Row Gov has 2 of the 3 lowest values.

### Drilling down education

5	Assoc	Gov	Private	Self-emp
	Assoc-acdm	39.91 (182)	40.87 (720)	45.49 (105)
	Assoc-voc	41.61 (169)	41.20 (993)	47.55 (145)
	Post-grad	Gov	Private	Self-emp
	Doctorate	46.53 (124)	49.05 (172)	47.22 (79)
	Masters	42.93 (567)	44.42 (863)	47.25 (197)
	Some-college	Gov	Private	Self-emp
	Some-college	38.38 (955)	38.73 (5016)	45.70 (704)
	University	Gov	Private	Self-emp
	Bachelors	41.56 (943)	42.71 (3455)	46.23 (646)
	Prof-school	48.40 (86)	47.96 (247)	47.78 (209)

2		Post-Secondary	Without Post-Secondary
	Gov	41.12	38.97
	Private	41.06	39.40
	Self-emp	46.39	44.84

### Summary for education

Act I  
(sl. 2,3)

*In this slide, we drill-down one level for all values of dimension work at level 0. For each cell we show both the Avg of Hrs and the number of tuples that correspond to it in parentheses. ...*  
Column Post-grad has 4 of the 6 highest values.  
Column Some-college has 4 of the 6 lowest values.

*In this graphic, we put the original request in context by comparing the value 'Post-Secondary' for education at level 3 with its sibling values. We calculate the Avg of Hrs while fixing education at level 4 to be equal to "ALL", and work at level 2 to be equal to "With-Pay". We highlight the reference cells with bold, the highest value with red and the lowest value with blue.*

Compared to its sibling we observe that in 3 out of 3 cases Post-Secondary has higher value than Without-Post-Secondary.

3		Assoc	Post-grad	Some-college	University
	With-Pay	41.62	44.91	39.41	43.44
	Without-pay	50.00	-	35.33	-

### Summary for work

### Drilling down work

4	Gov	Assoc	Post-grad	Some-college	University
	Federal-gov	41.15 (93)	43.86 (80)	40.31 (251)	43.38 (233)
	Local-gov	41.33 (171)	43.96 (362)	40.14 (385)	42.34 (499)
	State-gov	39.09 (87)	42.93 (249)	34.73 (319)	40.82 (297)
	Private	Assoc	Post-grad	Some-college	University
	Private	41.06 (1713)	45.19 (1035)	38.73 (5016)	43.06 (3702)
	Self-emp	Assoc	Post-grad	Some-college	University
	Self-emp-inc	48.68 (72)	53.05 (110)	49.31 (223)	49.91 (338)
	Self-emp-not-inc	45.88 (178)	43.39 (166)	44.03 (481)	44.44 (517)







## An example of a slide...

	Assoc	Post-grad	Some- college	University
Gov	40.73	43.58	38.38	42.14
Private	41.06	45.19	38.73	43.06
Self-emp	46.68	47.24	45.70	46.61

Here, you can see the answer of the original query. You have specified education to be equal to 'Post-Secondary', and work to be equal to 'With-Pay'. We report on Avg of work hours per week grouped by education at level 2. and work at level 1 .

You can observe the results in this table. We highlight the largest values with red and the lowest values with blue color.

Column Some-college has 2 of the 3 lowest values.

Row Self-emp has 3 of the 3 highest values.

Row Gov has 2 of the 3 lowest values.

# Cinecubes resources

Panos Vassiliadis

Dept. of Computer  
Science & Engineering  
Univ. Ioannina, Hellas

[pvassili@cs.uoi.gr](mailto:pvassili@cs.uoi.gr)

## Readings, Presentations and Demo

<http://www.cs.uoi.gr/~pvassili/projects/cinecubes/>

## Code

<https://github.com/DAI-NTI-NESS-Group/CinecubesPublic.git>

Dimitrios Gkesoulis, Panos Vassiliadis, Petros Manousis. CineCubes: Aiding data workers gain insights from OLAP queries. **Information Systems, Volume 53, October-November 2015**, Pages 60 - 86.

Dimitrios Gkesoulis, Panos Vassiliadis. CineCubes: Cubes as Movie Stars with Little Effort. **DOLAP 2013**, pp. 3 - 10, 28 October 2013, San Francisco, USA

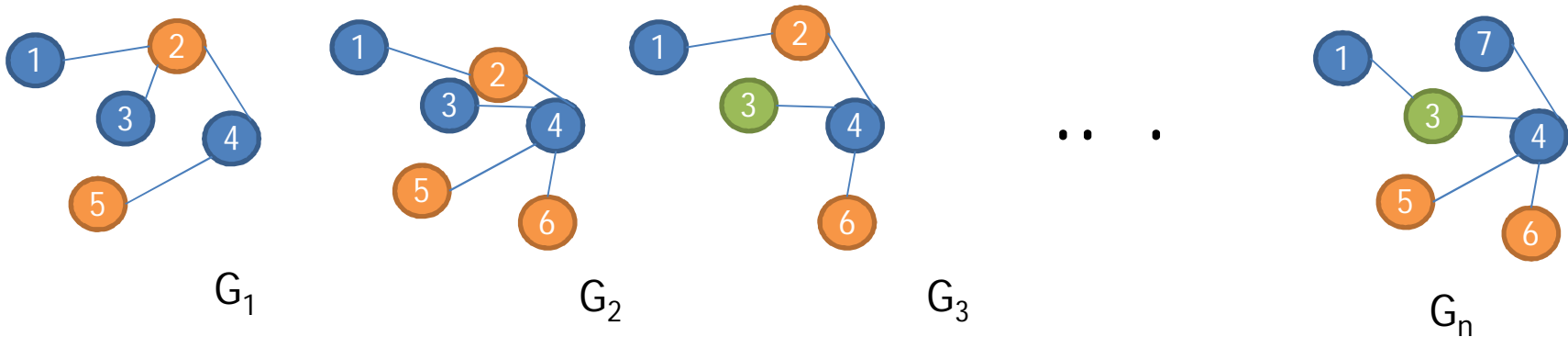
# Outline

- CineCubes: Aiding data workers gain insights from OLAP queries (Panos Vassiliadis, [pvassil@cs.uoi.gr](mailto:pvassil@cs.uoi.gr))
- Analysis of large social networks that evolve over time (Evaggelia Pitoura, [pitoura@cs.uoi.gr](mailto:pitoura@cs.uoi.gr))
- Extracting Knowledge from Social Networks and Their Content (Panayiotis Tsaparas, [tsap@cs.uoi.gr](mailto:tsap@cs.uoi.gr))
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# Time-Evolving Graphs

*Evolving graph*: A sequence of graph *snapshots*  $G_t$  at time instance  $t$



- § Storage issues
- § Indexing issues
- § Many novel historical graph queries that:
  - (a) have a time-range dimension (*when*),
  - (b) consider the durability of results (*how long/how often*), and
  - (c) capture time evolution (e.g., monitor).

# Pattern Matching

Given: **graph**  $G(V, E, L)$ ,  $L: V \rightarrow \Sigma^*$

**pattern**  $P(V_p, E_p, L_p)$

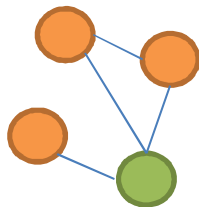
Find all subgraphs  $m = (V_m, E_m, L_m)$  of  $G$ , such that, there exists a *bijective function*  $f: V_p \rightarrow V_m$ :

$V_p \rightarrow V_m$ :

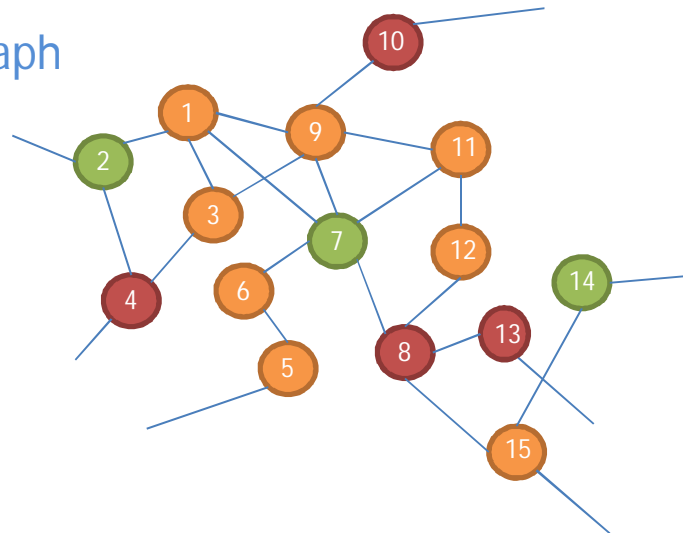
- for all  $u$  in  $V_p$ ,  $L_p(u) = L_m(f(u))$  and
- for each edge  $(u, v) \in E_p$ ,  $(f(u), f(v)) \in E_m$

Graph  $m$  is called a **match** of  $P$  in  $G$

**Pattern**



**Graph**



color - label

Identify the **most durable matches**: the matches that exist for the largest time interval, either collectively (i.e., in the largest number of graph snapshots), or continuously (i.e., in consecutive graph snapshots)

# Definition

Duration of a set of time intervals  $I$

§ *collective duration*: the number of time instants in  $I$

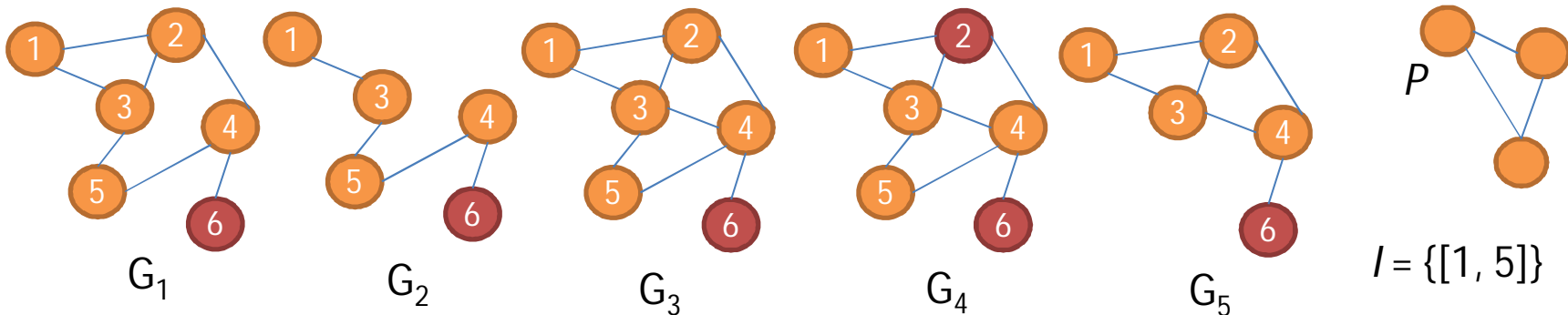
§ *continuous duration*: the duration of the longest time interval in  $I$

Example  $I = \{[1, 3], [5, 10], [12, 13]\}$  – Collective: 11, Continuous: 6

(Durable Graph Pattern Matching): Given an evolving graph  $G[i, j]$ , a graph pattern query  $P$  and a set  $I$  of time intervals:

- A *collective-time durable graph pattern query* finds the matches  $m$  such that  $\text{lifespan}(m) \cap I$  has the largest collective duration.
- A *continuous-time durable graph pattern query* finds the matches  $m$  such that  $\text{lifespan}(m) \cap I$  has the largest continuous duration.

time interval intersection/ lifespan: set of time intervals that an element (a node/edge/match, etc) exists



match [1 2 3] lifespan:  $\{[1, 1], [3, 3], [5, 5]\}$  top-collective with collective duration 3

match [3 4 5] lifespan  $\{[3, 4]\}$  top-continuous with continuous duration 2



# Durable Graph Pattern algorithm

A Filter-and-Verify algorithm based on three basic concepts :

1. *Version Graph* representation of the snapshot sequence (life-span annotated union graph of the sequence with an efficient in-memory representation using bit-arrays)
2. *Time Graph Indexes* used to filter candidate for matching nodes and refining them
3.  *$\vartheta$ -duration threshold* that dynamically estimates the minimum duration of a candidate match using the indexes

# Historical Reachability Queries

## Problem definition

Given nodes  $u$  and  $v$  and a set of time intervals  $I$  are  $u$  and  $v$  reachable in  $I$

- § *Disjunctive*, in at least one time instant in  $I$
- § *Conjunctive*, in all time instants in  $I$

## Key concepts

- § Find *Strongly-Connected-Components* (SCC) in each graph snapshot
- § Use bi-partite matching to *map* SCC at different snapshots effectively
- § Store reachability for SCC + extended *2HOP* index

# For more information

- K. Semertzidis, and E. Pitoura: *Durable Graph Pattern Queries on Historical Graphs*, ICDE 2016
- K. Semertzidis, E. Pitoura, K. Lillis: *TimeReach: Historical Reachability Queries on Evolving Graphs*, EDBT 2015

## Older papers

- K. Semertzidis, and E. Pitoura: *Time Traveling in Graphs using a Graph Database*, in Proc. of the 5<sup>th</sup> International Workshop on Querying Graph Structured Data (GraphQ 2016), in conjunction with the EDBT/ICDT 2016
- § G. Koloniari, E. Pitoura: Partial view selection for evolving social graphs. *GRADES 2013*
- § G. Koloniari, D. Souravlias, E. Pitoura: On Graph Deltas for Historical Queries. Workshop on Online Social Systems (*WOSS 2012*), in conjunction with the VLDB 2012



# Outline

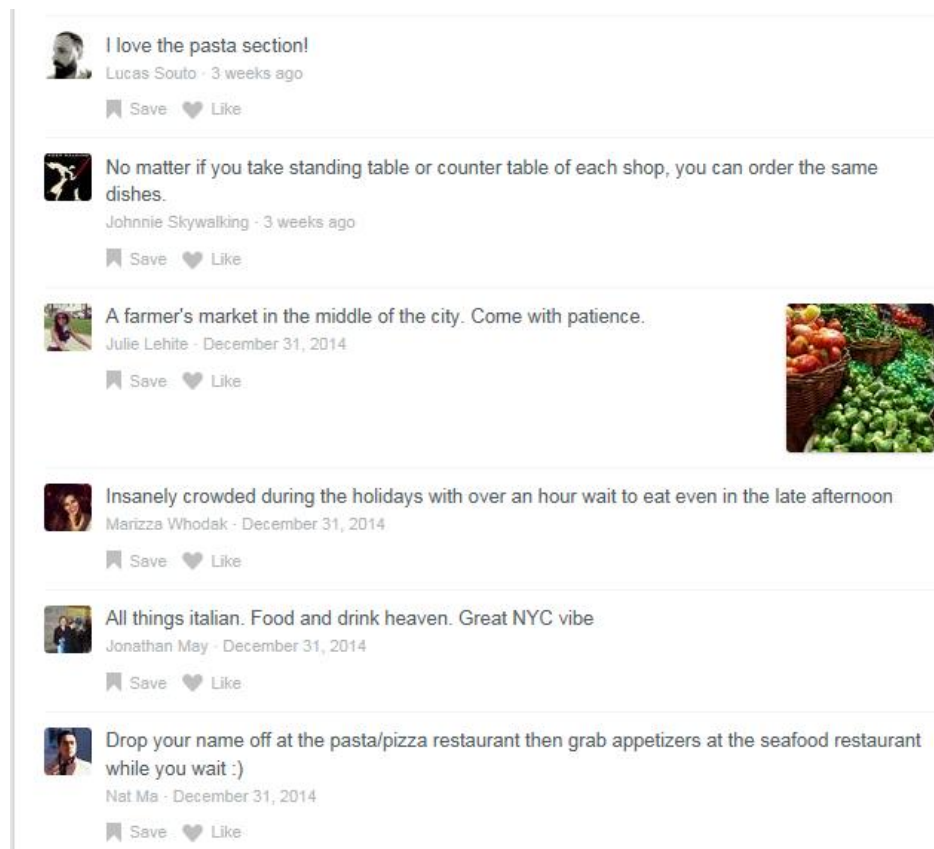
- CineCubes: Aiding data workers gain insights from OLAP queries (Panos Vassiliadis, [pvassil@cs.uoi.gr](mailto:pvassil@cs.uoi.gr))
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# Extracting Knowledge from Social Networks and Their Content

- Social Networks and Media produce huge volume of data with great commercial and scientific value
  - We need new algorithms for modeling and analyzing social networked data.
- Two areas of research
  - Mining user-generated reviews and micro-reviews
  - Understanding relationships and dynamic processes in networks.

# Mining Reviews and Micro-Reviews

- **Micro-reviews** (tips): A new type of User Generated Content!
  - Bite-size reviews (usually under 200 characters) commonly posted on social media or check-in services.

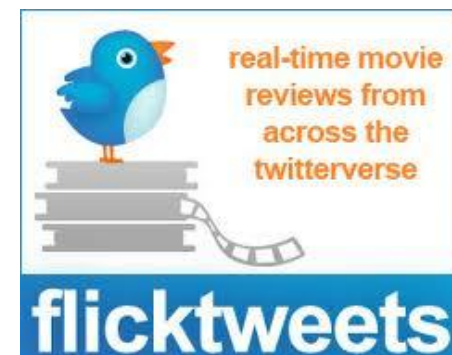


foursquare

facebook.



Facebook Places  
Who. What. When. And now where.



# Micro-Review Summarization

- Micro-reviews are highly informative but:
  - Short, repetitive, and in large volumes– tedious to go through them
  - Mostly consumed on mobile devices with limited screen space
    - hard to go through them
- We need a summary that:
  - Covers the salient points in the micro-reviews
  - Reasonable length
  - Flowing text, easy to read



# Reviews vs Micro-Reviews

## Reviews



off-site

elaborate and comprehensive

well-written,  
narrative/descriptive flow

## Micro-Reviews



on-site check-ins

concise and distilled

brusque and curt

Reviews and Micro-Reviews are complementary to each other

Goal: Use Reviews to summarize micro-review content



# Using reviews to summarize micro-reviews

- Two approaches:
  - Use micro-reviews to **select** a **small set of reviews** that capture the review content
    - Model the problem as a **coverage** problem where we want to **cover** the micro-reviews while not adding **redundant** content  
[T-S. Nguyen, H. W. Lauw, P. Tsaparas, *Review Selection Using Micro-reviews*. IEEE Transactions on Knowledge and Data Engineering, TKDE, 27(4), 1098-1111, 2015.]
  - **Synthesize** a **new review** from review snippets that **capture** the information in the micro-reviews in a **compact** form.
    - Use **MDL** to find a **compact** and **representative** subset of snippets  
[T-S. Nguyen, H. W. Lauw, P. Tsaparas, *Review Synthesis for Micro-Review Summarization*. ACM International Conference on Web Search and Data Mining (WSDM), 2015.]

# Dynamic Processes on Networks

- Diffusion phenomena on networks have been studied extensively as models for real-world phenomena
  - Viral marketing
  - Epidemic Spreads
  - Public opinion formation.
- An important and well studied problem in this area is the problem of diffusion maximization
  - Find a small set of nodes in the network that will be the initiators such that they maximize the spread of the diffusion on the underlying network
  - Greedy algorithms have provable approximation guarantees and work well in practice

[D. Kempe, J. Kleinberg, E. Tardos. *Maximizing the Spread of Influence through a Social Network*. Proc. 9th ACM SIGKDD Intl. Conf. on Knowledge Discovery and Data Mining, 2003.]

# Diffusion maximization

- Diffusion maximization over **dynamic networks**
  - When the network changes over time the Greedy algorithm no longer has provable guarantees
  - It is not only about **who** are the initiators, but also **when** they become active.  
[N. Gayraud, E. Pitoura, P. Tsaparas, *Diffusion Maximization on Evolving networks*, COSN 2015]
- Maximizing positive **opinions** on social networks
  - Opinions have continuous values as opposed to discrete adoption behavior.
  - **Opinion formation models** capture the process of opinion diffusion in a social network.
  - Find users to influence to change their opinion to positive so as to **maximize the overall positive opinion**.  
[A. Gionis, E. Terzi, P. Tsaparas. *Opinion Maximization in Social Networks*. SDM 2013]

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# Search and Analysis of spatially-enriched data

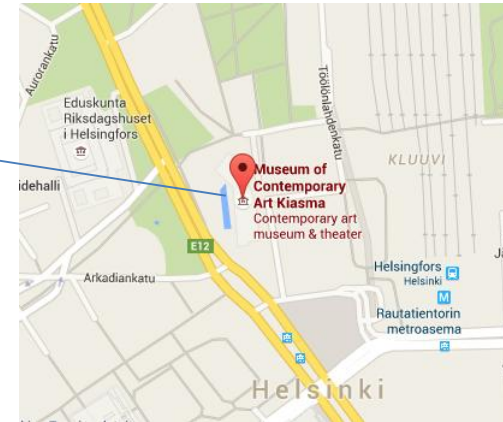
Location information in  
RDF knowledge bases

<i>subject</i>	<i>property</i>	<i>object</i>
Dresden	cityOf	Germany
Prague	cityOf	CzechRepublic
Leipzig	cityOf	Germany
Dresden	hosted	Wagner
Leipzig	hosted	Bach
Wagner	hasName	"Richard Wagner"
Wagner	performedIn	Leipzig
Dresden	hasGeometry	"POINT (...) "
Prague	hasGeometry	"POINT (...) "
Leipzig	hasGeometry	"POINT (...) "
...	...	...

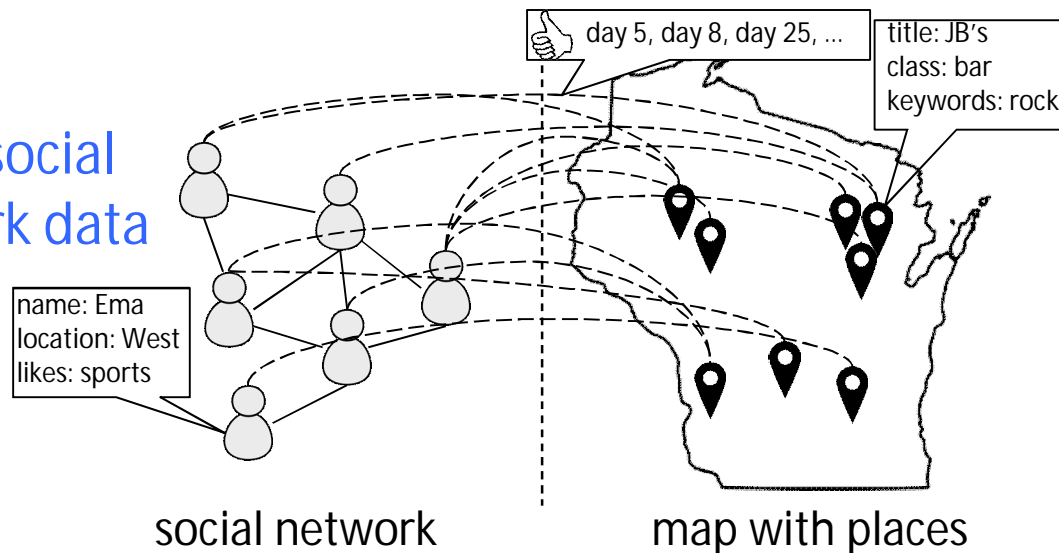
spatial  
entities

Geo-tagged documents or  
textually annotated POIs

"Contemporary  
art museum &  
cultural center  
with thematic  
annual  
exhibitions, a  
theater and  
events"



Geo-social  
network data



social network

map with places



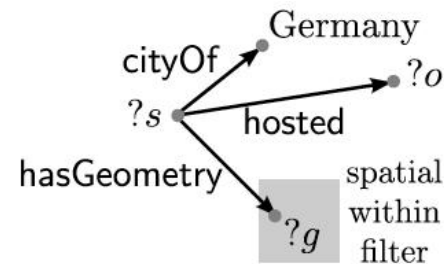
# Indexing Locations in large RDF data

1. *Encode locations of spatial entities* into their IDs

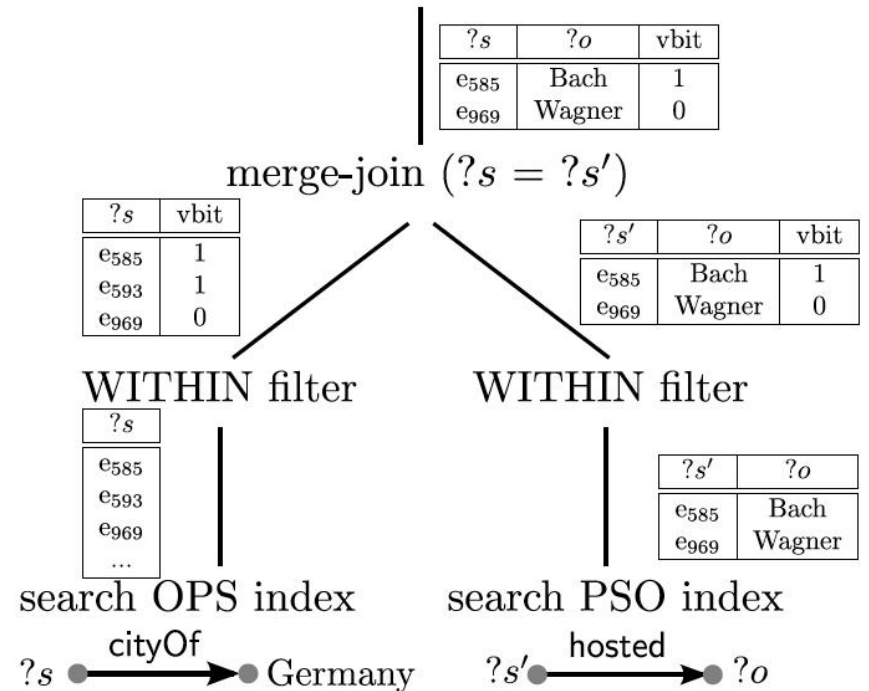
- i. Handle points/polygons
- ii. Handle spatial skew

2. *Extend RDF-3X* to support Spatial RDF queries

- i. On-the-fly spatial filters using entity IDs
- ii. Spatial join algorithm that operate on IDs
- iii. Extended query optimizer



spatial selection query



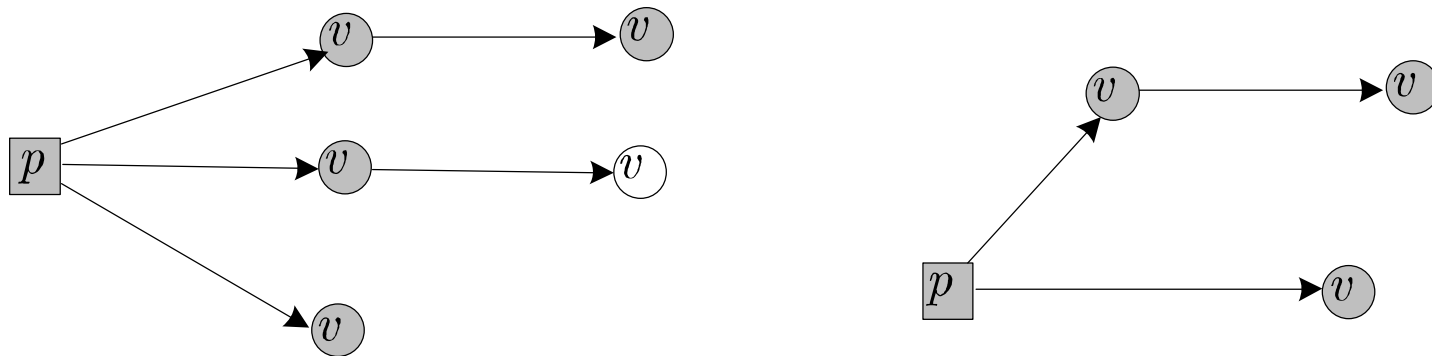
query evaluation plan

# Spatial keyword queries on RDF data

**Problem:** Find **places** in RDF data which are close to a query location and they are related to a set of given query keywords.

## Key concepts

- § Search for RDF subgraphs (1) **rooted at place entities** near query location (2) **containing query keywords**
- § Measure relevance to keywords by aggregated graph distance of keyword appearance to root
- § Spatial-first search, with the help of keyword reachability index / preprocessing



places related to {ancient,roman,catholic,history}

# Location Aware Keyword Query Suggestion

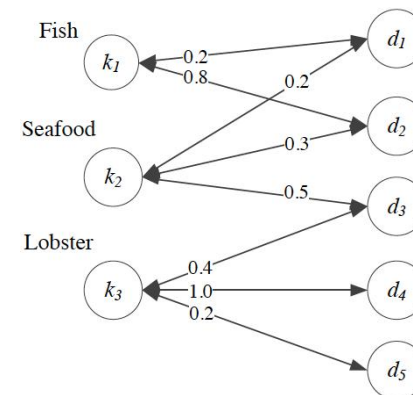
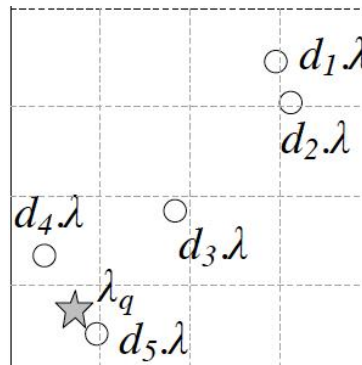
**Problem:** Suggest alternative keyword **queries** based on:

1. semantic relevance to original query and
2. proximity of results to query location

## Key concepts

- § A **Keyword-Document graph** connects past keyword queries to documents
- § Weights in the KD graph model query-document relevance (from click data)
- § Given a keyword query  $k_q$  and a location  $\lambda_q$ , edge weights are **adjusted**
- § Location-aware relevant queries to  $k_q$  are modeled by their **RWR distance** to  $k_q$
- § The graph is partitioned for more efficient RWR-based query suggestion

$d_1$	Fish and Seafood
$d_2$	Fish Seafood
$d_3$	Lobster Seafood
$d_4$	Lobster Restaurant
$d_5$	Lobster House
$k_1$	Fish
$k_2$	Seafood
$k_3$	Lobster

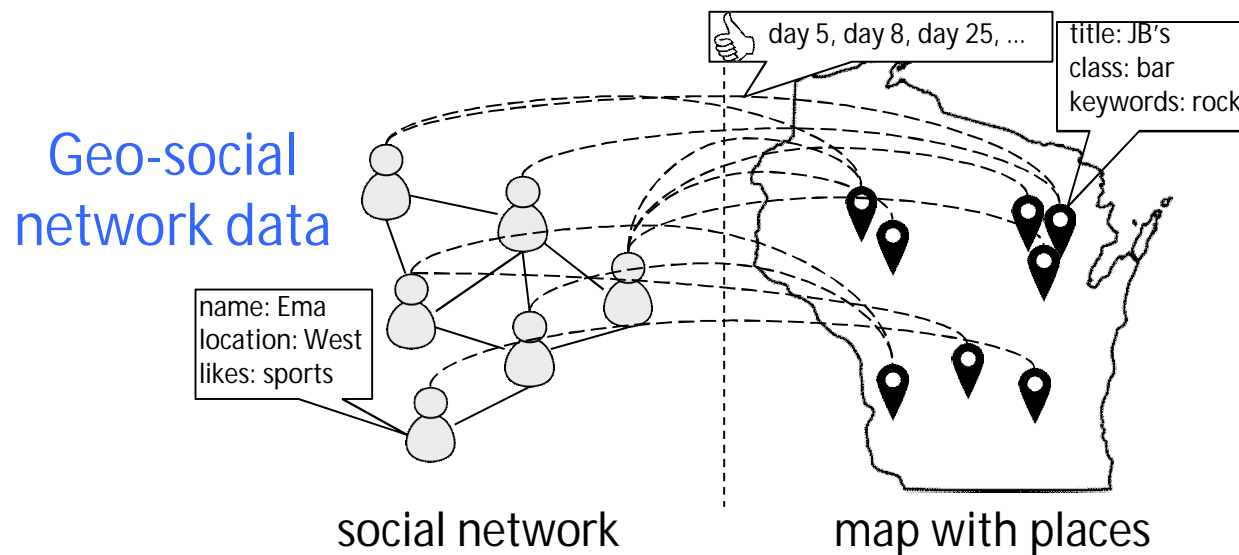


# Location Recommendation using Check-in Data

**Problem:** Recommend places to a geo-social network user

**Our recommender system considers**

- § **locations** of recommended venues (faraway venues are disregarded)
- § **similarity between check-in histories** of users (artificial weighted edges are added to the social graph connecting users with similar profiles)
- § **social relationships** between users (RWR graph distance is used to model similarities between users and then CF is applied)

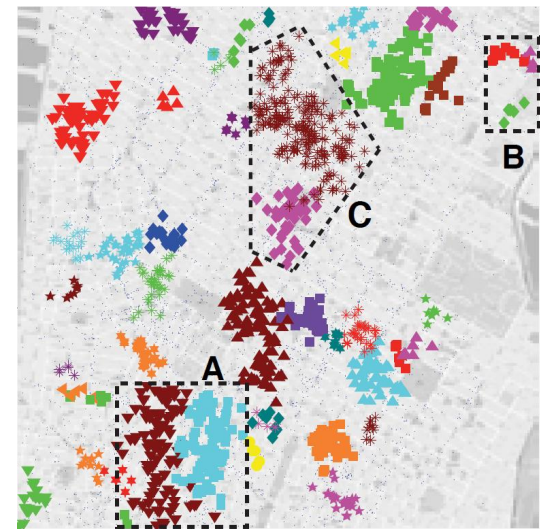
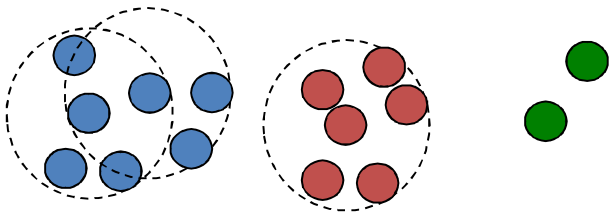


# Clustering Places in Geo-Social Networks

**Problem:** Consider the geo-social network when clustering places in an urban map

## Key concepts

- § **Social distance** between places  $D_S(p_i, p_j) = 1 - \frac{|CU_{ij}|}{|U_{p_i} \cup U_{p_j}|}$   
 $CU_{ij}$ : Contributing users for  $(p_i, p_j)$
- users who visited both places
  - users who visited one place AND have a friend who visited the other
- § **Geo-social distance** between places
- weighted sum of social and spatial distance
- § **Clustering** extends DBSCAN to use geo-social distance





# Joint Search by Social and Spatial Proximity

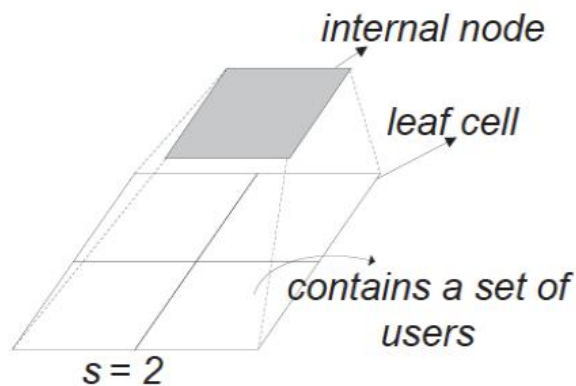
**Problem:** Find mobile users who are **geographically close** and **socially near** the target user

## Key concepts

§ a **hybrid ranking function**, which combines both distances

$$f(u_i, u_j) = \alpha \cdot E(u_i, u_j) + (1 - \alpha) \cdot D_S(u_i, u_j)$$

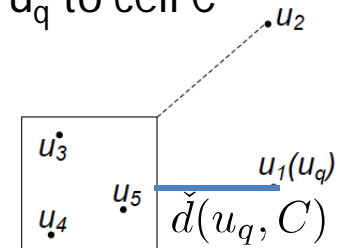
§ A multi-level spatial grid index with **social summaries** (aggregated distances to **landmarks** in the social graph)



$m_{ij}$  : distance between node  $v_i$  and  $j$ -th landmark  
 for a cell  $C$ :  $\hat{m}[j] = \max_{v_i \in C} m_{ij}$ ,  $\check{m}[j] = \min_{v_i \in C} m_{ij}$

define **lower distance bounds** from query user  $u_q$  to cell  $C$

$$\check{p}(u_q, C) = \max_{1 \leq j \leq M} \begin{cases} \check{m}[j] - m_{qj} & \text{if } m_{qj} < \check{m}[j] \\ m_{qj} - \hat{m}[j] & \text{if } m_{qj} > \hat{m}[j] \\ 0 & \text{otherwise} \end{cases}$$



**prioritize** examination of cells and their contents based on

$$MINF(u_q, C) = \alpha \cdot \check{p}(u_q, C) + (1 - \alpha) \cdot \check{d}(u_q, C)$$

K. Mouratidis, J. Li, Y. Tang, and N. Mamoulis, "Joint Search by Social and Spatial Proximity," IEEE TKDE + ICDE 2016 poster.

# More Info

[http://dmod.eu/data\\_web/](http://dmod.eu/data_web/)